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? t 24/ 7/ 1-2
               (Item 1 from file: 2)
DIALOG(R) File 2: INSPEC
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10562911
 Title: Bluetooth network - the ad hoc network concept
  Author(s): Suri, P.R.; Rani, S.
           Affiliation:
                              Dept. of Comput. Sci. & Applications, Kurukshetra
  Aut hor
Univ., Haryana, India
Conference Title:
                              Proceedings. IEEE Southeast Con 2007 (IEEE Cat.
No. 07CH37882)
                    p. 1 pp.
  Publisher: IEEE, Piscataway, NJ, USA
Publication Date: 2007 Country of Publication: USA
  ISBN: 1-4244-1028-2 Material Identity Number: XX2007-00845 U.S. Copyright Clearance Center Code: 1 4244 1028 2/2007/$25.00
  Conference Title: Proceedings. IEEE Southeast Con 2007
  Conference Date: 22-25 March 2007
                                                   Conference Location: Richmond, VA,
  Language: English
                             Document Type: Conference Paper (PA)
  Treatment: Practical (P)
                Ad hoc network is often local area network or other small area
  Abstract:
network formed by wireless devices. In Latin, ad hoc literally means "for this," further meaning "for this purpose only," and thus usually temporary. The area of ad hoc networking has gathered much research interests in the
past years. Bluetooth is one of the technologies that can be used for ad
     networking. The original idea of Bluetooth concept was that of cable
replacement between portable and/or fixed electronic device. According to
the specification, when two Bluetooth devices come into each other's communication range, one of them assumes the role of master of the
communication range, one of them assumes the role of master of the communication and the other becomes the slave. This simple "one hop"
network is called a piconet, and may include up to seven active slaves
connected to one master. As a matter of fact, there is no limit on the
maximum number of slaves connected to one master but only seven of them can
be active at time, others have to be in so called parked state. The
         unit of a piconet controls the traffic within the piconet by means
master
of polling the slaves according to any preferred algorithm e.g. Round Robin, which determines how the bandwidth capacity will be distributed among the slave units. The polling of slaves within a piconet results in
                 of the slaves in the master unit, which is referred to as scheduling . In this paper the usage of Bluetooth ad hoc
 schedul i ng
intra-piconet
net working in communication is elaborated. Instead of large-scale networks,
small-scale personal area networks are emerging in response to the
introduction of short-range radio technologies that is Bluetooth. (O Refs)
  Subfile: B C
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(Item 2 from file: 2) DIALOG(R) Filè 2: INSPEC (c) 2008 Institution of Electrical Engineers. All rts. reserv. INSPEC Abstract Number: B9403-6150M-031, C9403-5640-032 Title: End-to-end performance control for distributed real-time systems Author(s): Sholl, H.A.; Pia, P.J. Author`Affiliation: Connecticut Univ., Storrs, CT, USA Conference Title: Proceeding of the Twenty-Sixth Hawaii International Conference on System Sciences (Čat. No. 93TH0501-7) p. 463-72 vol. 2 Editor(s): Mudge, T.N.; Mlutinovic, V.; Hunter, L Publisher: IEEE, Los Alamitos, CA, USA Publication Date: 1993 Country Publication: USA 4 vol. (xvi +895+xi v+691+xi i+654+xv+889) pp. ISBN: 0 8186 3230 5 U. S. Copyright Clearance Center Code: 0-8186-1060-3425/93/\$03.00 Conference Sponsor: ACM, IEEE Conference Date: 5-8 Jan. 1993 Conference Location: Wailea, HI, USA Document Type: Conference Paper (PA) Language: English Treatment: Practical (P) Abstract: A best effort scheduling algorithm is used to reduce the probability of exceeding a deadline-related target time for each job class allocated to a **distributed** real-time system A dynamic **scheduler** is used as a point of control on each node of a pipeline-structured **distributed** system The authors contrast a local clustering **scheduling** algorithm which attempts to meet end-to-end timing requirements in isolation to a global clustering scheduling algorithm which attempts to meet end-to-end timing requirements by utilizing information about the state of other nodes in a job's execution path. The approach incorporates both job-class-based loss functions and feedback of remaining time estimates. A simulation study has shown that the global algorithm can provide an even distribution of processing power over all nodes and job classes while maintaining system stability. (24 Refs) Subfile: B C